# **Optimized Ring Routing Protocol for Efficient Data Transmission**

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Abstract - In Wireless Sensor Network (WSN) energy eff iciency is the main issue. Because the nodes near the mobil e sink are reduce faster than the other nodes. The reason is all the data requests are transferred towards the same no de near to the mobile sink. To overcome this problem there are many approaches are already established which aims to decrease the overload of advertising the sink's position to the network by implementing a virtual hierarchy of nod es which imposes different dynamic roles on the sensor no des. The existing Grid-Based Energy-Efficient Routing fro m Multiple Sources to Multiple Mobile Sinks (GBEER) appr oach spreads data declarations straight along the shared grid while data requests are propagated precipitously, ens uring that these packets intersect at a crossing point. The position of the sink is then delivered to the source node, an d data is delivered directly to the sink. However, constructi on of the grid is non-trivial and some other techniques are also exist non-hierarchical approaches which utilize mech anisms like overflowing, overhearing or exploit geometric properties to advertise the sink's position to the source no des. To enhance the efficient way of routing to exhaust the energy used by WSN nodes, introducing Ring Routing is a hierarchical mobile sink routing protocol based on a virtu al ring structure which is designed to be easily accessible a nd easily reconfigurable. In this first ring structure is const ructed based on the location of the mobile sink. The results show that Ring Routing indeed is an energy-efficient prot ocol which extends the network lifetime.

*Key Words: Wireless Sensor Network, mobile sink, ener gy efficiency, routing protocol.* 

# **1. INTRODUCTION**

In Wireless Sensor Network there are many approaches t o the problem of routing in WSNs with mobile sinks. The most important and the most widely adopted one being are the hierarchical routing[7] protocols. In a typical wir eless sensor network, the energy of the nodes near to the sink reduces faster than the other nodes due to the data transmission using the same path to the sink. In wireless sensor networks (WSNs), energy efficiency[2] is conside red to be a crucial issue due to the limited battery capacit y of the sensor nodes. This problem generally affects the trustworthy nodes in the network. It may also deplete th e whole wireless sensor network within a short time. He nce this is important to overcome the energy drain of bat tery reduction.

To overcome the Wireless Sensor Network problem many approaches are already established. This may inclu de the routing protocols[7] like flat routing and hierarchi cal routing. Flat routing protocol is only useful for small s cale systems but not for large scale systems. Flat routing protocols include Sensor Protocol for Information via Na vigation(SPIN), Directed Diffusion, Rumor routing, Mini mum Cost Forwarding Algorithm(MCFA), Gradient-Base d Routing, etc. Hierarchical routing protocols include Lo w Energy Adaptive Clustering Hierarchy (LEACH), Thres hold-sensitive Energy-efficient protocols (TEEN), Small Minimum Energy Communication Network (MECN), Sen sors Aggregates Routing, etc.

In this paper we focus on the hierarchy routing prot ocol. The main reason is that the hierarchical routing is u sed in large scale systems. Hierarchical routing protocol structure includes rectangular, cluster, quad tree, line, ra il and ring. These structures are having different kind of protocols and benefits.

# **2. SYSTEM MODEL**

In this section we consider the existing system design an d the proposed system.

# 2.1 Existing System

The existing design approaches are Sink Trail, Line Base d Data Dissemination(LBDD)[10], Two-Tier Data Dissem ination (TTDD)[13], Grid-Based Energy-Efficient Routing (GBEER)[14].

#### Sink Trail:

Sink Trail establishes a logical coordinate system for rou ting and forwarding data packets, making it suitable for diverse application scenarios. This technique aim to decr ease the load of advertising the sink's position to the net work by establishing a virtual hierarchy of nodes which i mposes different dynamic roles on the sensor nodes.

#### Line Based Data Dissemination (LBDD):

It defines a vertical strip of nodes horizontally centered o n the area of development. The nodes on this strip are re ferred to as in-line nodes. Sensor data are sent to the line and then the first in-line node encountered stores the da ta. The nodes can access the line by straight forward mec hanism.

#### **Two-Tier Data Dissemination (TTDD):**

It is a virtual grid based approach where each source nod e with sensor data proactively constructs a rectangular g rid around itself and becomes a crossing point of this gri d. Whenever sinks require data, they query the network by local flooding within a grid cell and these queries are relayed to the source node. Data is then forwarded to the sink using the reverse of the path taken by the data requ est.

#### Grid-Based Energy-Efficient Routing (GBEER):

Data announcements are propagated horizontally along t he shared grid while data requests are propagated vertic ally, ensuring that these packets intersect at a crossing p oint. The position of the sink is then delivered to the sour ce node, and data is delivered directly to the sink. Grid-b ased protocols are advantageous for the easy accessibilit y of the grid structure.

#### **Problem Identified:**

- TTDD suffers from the high overhead of constructing a separate grid for each source node especially in applications where numerous sensor nodes generate data.
- GBEER eliminates the high overhead of constructing separate grids for each source, the nodes residing on the grid are likely to become hotspots and deplete their energy quicker than other nodes.
- GBEER requires more cost for implementing the mobile sink.

#### 2.2 Proposed System

Ring Routing is a hierarchical routing protocol for large s cale WSNs organized outside with static sensor nodes an

d a mobile sink. Ring Routing protocol first construct a vi rtual ring structure that allows the fresh sink position to be easily delivered to the ring and sink nodes to obtain t he sink position from the ring with minimal overhead wh enever needed.



#### Fig-1: Ring Routing System Design

The ring structure can be easily changed. The ring no des are able to change roles with regular nodes by a dire ct and well-organized mechanism, thus modifying the ho tspot difficult.

#### **Benefits:**

- Dynamic changing of ring structure will provide the efficient accuracy.
- Moderate cost for implementing mobile sink.
- Most secure and efficient communication.

# **3. DESIGN CONSTRUCTION**

This Section consists of the following module design to fo rm the ring routing protocol. These are to be explained i n this section.

#### **3.1 Ring Composition**

Ring Composition is dependent on the location informati on of the nodes, which is known to contain some incorre ctness based on the developed technology. Monte-Carlo analysis to determine the successful ring construction lik elihood under varying degrees of localization error. Net work Center set the Radius for forming the closed loop c ontaining several nodes. These nodes are called "Ring No des".



Fig-2: Ring Composition

Ring nodes are selected based on the distance of the node from the network authority. Monte-Carlo analysis i s used for successful ring construction.

# **3.2 Announcement of Sink Position**

Initially, the sink selects the contiguous node as it's AN, a nd advertises an AN Selection (ANS) packet. Before the si nk leaves the consultation range of the AN, it selects a ne w AN and informs the old AN of the position. Since now t he old AN knows about the new AN, it can relay any data which is designed for it to the new AN.

The current AN relays data packets directly to the si nk. After a ring node receives an ANPI packet, it shares th is information by sending an AN Position Information Sh are (ANPIS) packet to its clockwise and counter-clockwis e ring neighbors.



**Fig-3: Announcement of Sink Position** 

Each ring node receiving an ANPIS packet relays it to the neighbor ring node in the respective direction until t he two ANPIS packets sent in the clockwise and counterc lockwise directions arrive at the same ring node.

### **3.3 Anchor Node Position Information Nego** tiation

Upon selection of a new AN, it sends an AN Position Infor mation (ANPI) packet in the direction of the ring. If the A N is exterior to the ring, it sends the ANPI packet to the n etwork center, and if it is internal to the ring, it sends it d ata towards a point which be inherent in the opposite dir ection of the network center. The source node sends an AN Position Information Request (ANPIREQ) packet in th e direction of the ring. The ring node receiving the ANPI REQ packet generates an AN Position Information Respo nse (ANPIRESP) packet which include the current AN's p osition and sends the data to the source node.



#### Fig-4: Anchor Node Position Information Negotiation

# 3.4 Data Delivery

In this module source node receives the response from t he ANPIREQ. If the source node get the response from th e anchor node it knows the position of the AN and can no w send its message directly to it by geographic forwardin g. If data reaches an old AN, that means that the AN has a lready changed by the time data has arrived at the destin ed AN, the follow-up mechanism is used to transmit data to the current AN.





#### 4. PERFORMANCE EVALUATION

Ring Routing has the best performance in all cases. LBDD performs better than Railroad for sink speed . LBDD's av erage energy consumption tends to increase monotonica lly. LBDD employs broadcasts along the line structure to share AN position information, thus increased rate of AN changes caused by increased sink speed leads to increase d number of broadcast and thus elevated energy consum ptions. Railroad limits the number of broadcasts along th e rail by constructing localized stations in which the broa dcasts are confined. The AN position information is shar ed by unicasts along the rail until a station constructed b y broadcasts is reached. Therefore, Railroad performs be tter than LBDD for faster sink speed values.



# Chart-1: Delay breakdown of Ring Routing data deli very compared to LBDD.

In order to investigate the delay cost of the ANPI req uest/response mechanism employed by Ring Routing, w e provide the delay breakdown of the data delivery proc ess in Chart-1. The total delay for data deliveries are bro ken down into two components. The ANPI request/respo nse delay per data component is the time until a respons e to the ANPI request is received by a source node. The s econd component is the actual data dissemination delay of the path from the source to the sink. The two componen nts of Ring Routing's data delivery delays are compared with LBDD's average reporting delays. The delay cost of the request/ response mechanism is apparent. The actua l data dissemination delays of Ring Routing is lower tha n LBDD's reporting delays.

Even though, discarding the request/response mech anism of Ring Routing and employing a direct data sending approach similar to LBDD would enhance the de lay performance of Ring Routing, the energy advantages of the request/response mechanism outweighs such a de lay reduction benefit.

Since the delay values of Ring Routing close to LBDD 's, and they are in reasonable limits to support time-sens

itive applications, the energy consumption performance i s favored.

Even though the average energy consumption metric provides an insight to the projected longevity of the WS N operation, the lifetime metric is the clearer indicator. We initially define the lifetime of a WSN to be the time u ntil the first node in the network dies.



Chart-2 shows the lifetimes of Ring Routing, LBDD a nd Railroad for varying sink speed values. LBDD, as expe cted, has the worst lifetime performance among the thre e protocols due to the decreased degree of the energy co nsumption uniformity caused by all the data traffic being handled by the line nodes. The lifetimes of Railroad are c lose to but slightly worse than Ring Routing. Railroad em ploys a request/response mechanism similar to Ring Ro uting; however, broadcasts along the second- tier structu re are not entirely avoided, thus enabling Ring Routing t o stand out in terms of the network lifetimes. The static s ink case is also provided to show the huge advantage of mobile sinks in terms of lifetime elongation. The definitio n used for the network lifetime is reasonable but not eno ugh to accurately assess the longevity of the WSN operati on. The death of the first node might disrupt the topolog y and cause disconnectedness of some portions of the ne twork, depending on the criticality of the dying node's p osition. However, these problems are implicitly mitigate d by the usage of a mobile sink, since the sink is expected to eventually visit the disconnected areas in the network , thus providing stronger reliability and data delivery per formance.

In Chart-3, the distribution of the node death times u ntil 25 percent (50 nodes) of the 200 node network dies i s provided. The death of 25 percent of the network, whic h is a significant portion, hinders the WSN operation gre atly, in terms of both the imminent topological disruptio ns and the decrease of the sensor field's coverage.



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#### **Chart-3: Death times**

The constant sink speed value of 3 km/h is used for t his set of simulations since the lifetime performances of t he three protocols are observed to be the closest in the 3 km/h sink mobility case. Fig.3 emphasizes the strength o f Ring Routing's lifetime performance. The deaths of the f irst 50 nodes in Ring Routing occur later than LBDD and Railroad on average. Moreover, for Railroad, the nodes di e quicker than Ring Routing even though the first node d eath times are similar. LBDD's first node death times are much lower than Ring Routing and Railroad; however, th e first 50 nodes death times distribution shows that LBD D's rate of node deaths decrease as more nodes die. LBD D performs better than Railroad after the death of 12 pe rcent of the network, due to the later node death times b eyond this instance. LBDD nearly catches up with Ring R outing as the 25 percent of the network dies; however, e valuating beyond this value is irrational since the succes sful operation of the WSN is very likely to be already inte rrupted. The sink speed is not the only parameter affecti ng the WSN performance. The network size in terms of t he number of deployed sensor nodes also affects the WS N performance significantly since the density of the netw ork and the total traffic loads depend on the network siz e.

# **5. CONCLUSION**

Ring Routing relies on minimal amount of broadcasts. Th erefore, it is applicable to be used for sensors utilizing as ynchronous low-power MAC protocols designed for WSN s. Ring Routing does not have any MAC layer requiremen ts except the support for broadcasts. It can operate with any energy-aware, duty cycling MAC protocol (synchroni zed or asynchronized). Ring Routing is suitable for both event-driven and periodic data reporting applications. It is not query based so that data are disseminated reliably as they are generated. Ring Routing provides fast data d elivery due to the quick accessibility of the proposed ring structure, which allows the protocol to be used for time sensitive applications. No information about the motion of the sink is required for Ring Routing to operate. It doe s not rely on predicting the sink's trajectory, and is suita ble for the random sink mobility scenarios. Hence this is the optimal solution for avoiding the sink node energy de pletion attack.

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